

### Lecture 9: Generalised Algebraic Data Types

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## **Exam Information**

Time: 8am-12pm AEST on Monday Aug 14.Length: 2 hours. Make sure you start before 10AM.Where: online. Link will appear on the course website.

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Material: all material that was presented in the course, including in lectures, practicals, exercise sets or quizzes (except where we explicitly told you that the material was not examinable).

Format: There will be quiz-style questions about design. There will be theory questions. We may ask you to write code and proofs, but no long-form software implementation.

Sample exam will be released on the course website shortly.

Generalized Algebraic Data Types (GADTs) is an extension to Haskell that, among other things, allows data types to be specified by writing the types of their constructors:

```
data Answer = Yes | No
-- is the same as
data Answer :: * where
Yes :: Answer
No :: Answer
```

# GADTS

We will need to use two new language extensions to declare them.

## Aside: Sum Types

```
data Parity = Even | Odd
data Polarity = Positive | Zero | Negative
data Sum :: * -> * -> * where
L :: a -> Sum a b
R :: b -> Sum a b
```



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Do we see why they are called sum types?

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Do we see why they are called product types? NB: here we count *non-bottom* elements; e.g. undefined doesn't count.

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## **Sized lists**

We can use GADTs+phantom types to encode the length of a list in its type:

```
data Size = Z | S Size
```

```
data Vec :: * -> Size -> * where
Nil :: Vec a Z
Cons :: a -> Vec a n -> Vec a (S n)
```

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- Nil always has length 0 (Z)
- Cons x xs is one longer than xs (S n)

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#### Observation

This subsumes the distinct types for empty and non-empty lists we've seen previously.

Look at the type of the map function for Vec:

```
mapV :: (a -> b) -> Vec a n -> Vec b n
```

```
mapV f Nil = Nil
```

mapV f (Cons x xs) = Cons (f x) (mapV f xs)

It says that if the input has length n, then so does the output. So the property that mapV preserves length is enforced by the type system!

Think about all the inductive proofs we don't have to write.

GADTs are one of the most powerful static assurance tools available in Haskell. But:

- It can be difficult to convince the Haskell type checker that your code is correct, even when it is.
- Type-level encodings can make types more verbose and programs harder to understand.
- Too detailed types can make type-checking very slow, hindering xsproductivity.

### Be pragmatic!

Use type-based encodings when the assurance advantages outweigh the potential disadvantages. The typical use case is to eliminate partial functions from our code base

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## That's all folks!

Thanks for taking the course.

Don't forget to take the myExperience survey.